

Extremal Problems in Eulerian Digraphs

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Graphs and digraphs (directed graphs) behave quite differently

- ▶ any graph G with n vertices and m edges has a subgraph with minimal degree $\geq m/n$, implying
- ▶ G has a cycle of length $\geq m/n + 1$,
- ▶ however, those facts are trivially false for general digraph, even for dense digraph, i.e:
 - ▶ transitive tournament,
 - ▶ complete bipartite digraph $A \cup B$ with all edges from A to B .
- ▶ Consider restricted family of digraphs

► **Eulerian digraph:**

- out-degree $d^+(u) =$ in-degree $d^-(u)$ for any vertex u .
- No loop and parallel edges; may have 2-cycles.

► why do we consider Eulerian digraph?

a natural candidate; "easy" to evaluate eigenvalues; it has good properties:

- (1) for any cut (A, B) , number of edges from A to B = number of edges from B to A :

$$e(A) + e(A, B) = \sum_{u \in A} d^+(u) = \sum_{u \in A} d^-(u) = e(A) + e(B, A)$$

- (2) the edge set can be decomposed into edge-disjoint cycles,

Some interesting problems

- ▶ Conjecture(Bollobás and Scott, Dean): for Eulerian digraphs, there exists a cycle decomposition with $O(n)$ cycles ?
 - ▶ Even not known for graph!
- ▶ Conjecture (Erdős, Goodman and Pósa, 1966): for even graphs, there exists a cycle decomposition with $O(n)$ cycles ?
 - ▶ easy to get $O(n \log n)$:

greedily pick and remove a cycle of length $\geq m/n$ in the current graph; keep doing (say t steps) until

$$m(1 - \frac{1}{n})^t = n \quad \Rightarrow \quad t \leq n \log m/n \leq n \log n.$$

- ▶ Conjecture(Bollobás and Scott, 1996): Any Eulerian digraph with m edges and n vertices has a cycle of length $\Omega(m/n)$.
 - ▶ this motivates our work

more natural parameters

- ▶ **girth** $g(G)$ = the length of the shortest cycle in G
- ▶ **minimal feedback arc set** = the smallest set of edges whose removal results in an acyclic digraph; denote its cardinality as $\beta(G)$.
 - ▶ important in combinatorial optimization and with many applications to other fields.
 - ▶ computing $\beta(G)$ is NP-hard even for tournaments.
- ▶ We are also interested to find sub-digraph whose minimal degree is large.
- ▶ For graphs, general bounds about those parameters are known.
 - ▶ DFS \Rightarrow long cycle of length $\Omega(m/n)$
 - ▶ Moore bound $\Rightarrow g(G) = O(\log n)$ (if $m \geq 3n$)
 - ▶ $\beta(G) = m - (n - 1)$ if G is connected

Our results

Every Eulerian digraph G with n vertices and m edges has

- ▶ **Thm 1:** $\beta(G) \geq \frac{m^2}{2n^2} + \frac{m}{2n}$. Optimal for any m, n with $m = kn$.

⇓

- ▶ **Thm 2:** girth $g(G) \leq 6n^2/m$. Tight up to a constant.

⇓

- ▶ **Thm 3:** an Eulerian sub-digraph with minimal degree $\geq m^2/24n^3$. Tight up to a constant.

⇓

- ▶ **Thm 4:** a cycle of length $\geq \max\{m^2/24n^3, \sqrt{m/n}\}$.

- ▶ When $m = cn^2$, this gives a cycle of length linear in n .
- ▶ When graph is sparse, $\sqrt{m/n}$ is better.

Thm 1: Minimal feedback arc set $\beta(G) \geq \frac{m^2}{2n^2} + \frac{m}{2n}$

- ▶ Extremal examples: there exists a Cayley digraph G for any m, n with $m = kn$ such that $\beta(G) = \frac{m^2}{2n^2} + \frac{m}{2n}$.
- ▶ A weaker bound what we are going to prove: $\beta(G) \geq \frac{m^2}{8n^2}$.
- ▶ consider a linear order of vertices, say v_1, v_2, \dots, v_n . Arc (v_i, v_j) is **backward** if $i > j$; otherwise **forward**.
- ▶ it suffices to prove: for arbitrary order of vertices, the number of backward arcs $\geq \frac{m}{4k}$, where $k = 2n^2/m$.
- ▶ A careful analysis will lead to the optimal bound: we need to count the “contribution” of an arc whose two endpoint are far. Complicate.

From $\beta(G) \geq \frac{m^2}{2n^2} + \frac{m}{2n}$ to Thm 2: girth $g(G) \leq 6n^2/m$

- ▶ Need a Theorem of Fox-Keevash-Sudakov: If a digraph G satisfies $\beta(G) > 18n^2/r^2$ for $r > 11$, then $g(G) \leq r$.
- ▶ If there exists a 2-cycle, then $g(G) = 2 \leq 6n^2/m$.
- ▶ Assume G has no 2-cycle, then $6n^2/m \geq 12 > 11$, then

$$\beta(G) > \frac{m^2}{2n^2} = \frac{18n^2}{(6n^2/m)^2} \quad \Rightarrow \quad g(G) \leq 6n^2/m.$$

- ▶ Tight up to constant and the optimal constant seems to be 1: again the Cayley digraphs.
 - ▶ Caccetta-Häggkvist Conjecture; when $r = n/3$ and Eulerian

Thm3: an Eulerian subgraph with large minimal degree $\geq m^2/24n^3$

- ▶ Thm 2 says that Eulerian digraph with at least $m/2$ edges has a cycle of length $\leq 12n^2/m$.
- ▶ at least $m/2$ edges can be decomposed into t edge-disjoint cycles of each length $\leq 12n^2/m$, thus $t \geq m^2/24n^2$.
- ▶ Define the union of cycles as H , which is an Eulerian digraph.
- ▶ If there is a vertex u whose degree $< m^2/24n^3$, then we delete u and all short cycles through u which we choose.
- ▶ When stop, we get a sub-digraph with such property.

Extremal digraphs in which the minimal degree of any Eulerian sub-digraph is $O(m^2/n^3)$.

- ▶ $V(H(s, t)) = U_1 \cup \dots \cup U_s \cup V_1 \cup \dots \cup V_t$ and $|U_i| = |V_j| = s$.

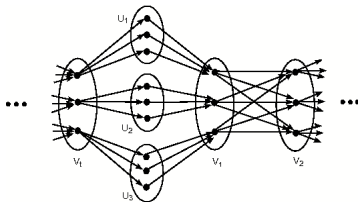


Figure: $H(s, t)$ with $s = 3$

- ▶ Any cycle in $H(s, t)$ must pass through some vertex in U_i , whose degree is 1.
- ▶ Our extremal digraphs: δ -blowup of $H(s, t)$, which has $n = s(s + t)\delta, m = s^2(t + 1)\delta^2$.
- ▶ Taking $t = 2s$, then $m^2/n^3 \geq 4\delta/27 \Rightarrow \delta \leq O(m^2/n^3)$.

Thm4: cycle of length $\geq \max\{m^2/24n^3, \sqrt{m/n}\}$

- ▶ Lemma: any digraph with minimal out-degree k has a cycle of length $\geq k + 1$.

\Downarrow

- ▶ a cycle of length $\geq m^2/24n^3$
- ▶ An inductive argument shows the existence of cycle of length $\geq \sqrt{m/n}$.
 - ▶ By lemma, we assume vertex u has out-degree $t \leq \sqrt{m/n}$, then we delete u and all cycles C_1, \dots, C_t passing through u . Also assume $|C_i| < \sqrt{m/n}$.
 - ▶ The remaining graph has $n - 1$ vertices and m' edges:

$$m' \geq m - t\sqrt{m/n} \geq m(1 - \frac{1}{n})$$

- ▶ By induction, there exists a cycle of length at least $\sqrt{m'/(n-1)} \geq \sqrt{m/n}$.

Open problems

- ▶ **Bollobás-Scott:** Should have a cycle of length $\Omega(m/n)$.
 - ▶ DFS fails
- ▶ **Caccetta-Häggkvist:** every digraph with n vertices and minimal out-degree at least r should contain a cycle of length $\leq \lceil n/r \rceil$.
 - ▶ Most interesting case: $r = n/3$, which is open even for Eulerian digraphs!
- ▶ **Seymour:** For Eulerian digraphs,

$$\sum_{i \in V} d^{++}(i) \geq \sum_{i \in V} d^+(i) = m.$$

Happy Birthday! Robin!